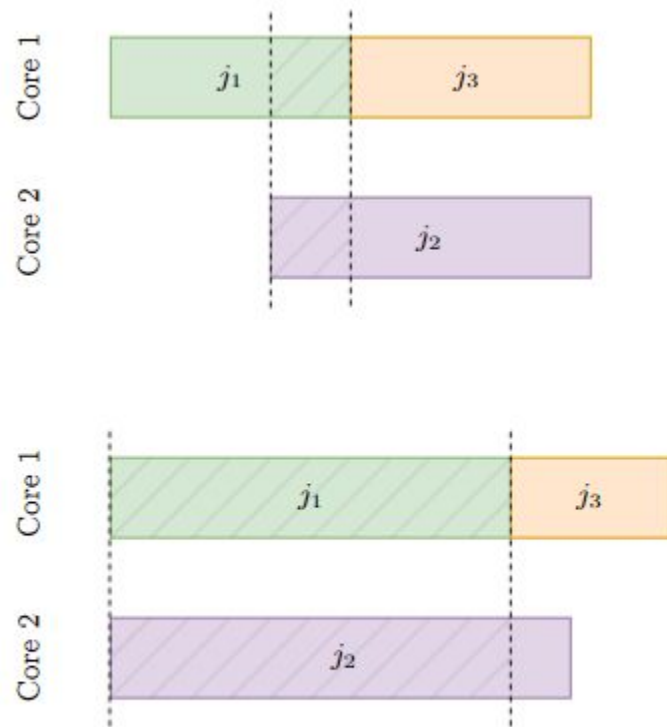


# On Different Methods for Automated MILP Solver Configuration

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# Base task



# Base task

$$C_{\max} = \sum_{k \in K} \sum_{c \in C} t_{kc} \rightarrow \min, \quad (1)$$

$$t_{kc} \geq 0, \quad k \in K, c \in C, \quad (2)$$

$$t_{kc} \leq x_{kc} T_{\max}, \quad k \in K, c \in C, \quad (3)$$

$$\sum_{c \in C} x_{kc} = 1, \quad k \in K, \quad (4)$$

$$\sum_{k \in K} \sum_{c \in C} t_{kc} v_{jc} = s_j, \quad j \in J, \quad (5)$$

$$\sum_{c \in C} x_{kc} q_{jc} - \sum_{c \in C} x_{k-1,c} q_{jc} \leq y_{jk}, \quad j \in J, k \in K, \quad (6)$$

$$\sum_{k \in K} y_{jk} = 1, \quad j \in J, \quad (7)$$

$$\sum_{k \in K} (k+1)x_{kc_1} \leq (1 - \sum_{k \in K} x_{kc_1})e + \sum_{k \in K} kx_{kc_2}, \quad c_1, c_2 \in C, a_{c_1 c_2} > 0, \quad (8)$$

$$x_{kc} \in \{0, 1\}, y_{jk} \in \{0, 1\}, \quad j \in J, k \in K, c \in C. \quad (9)$$

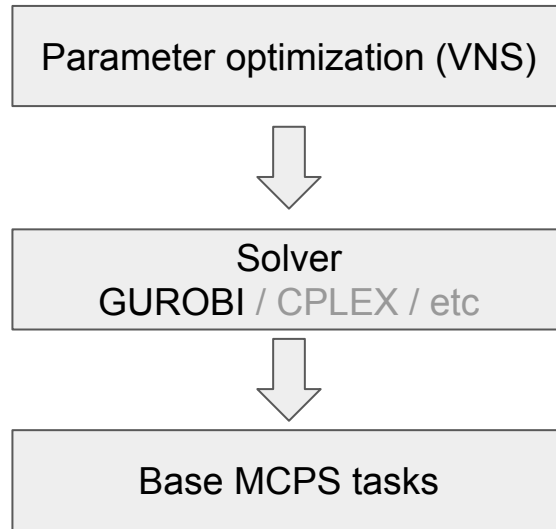
# Base task

1.  $q_{jc} = 1$  when and only when job  $j$  is included in configuration  $c$ , 0 else.
2.  $a_{c_1c_2} = 1$  when and only when configuration  $c_2$  must be performed after configuration  $c_1$ , 0 else.
3.  $T_{\max}$  — upper bound of configuration runtime in event point.

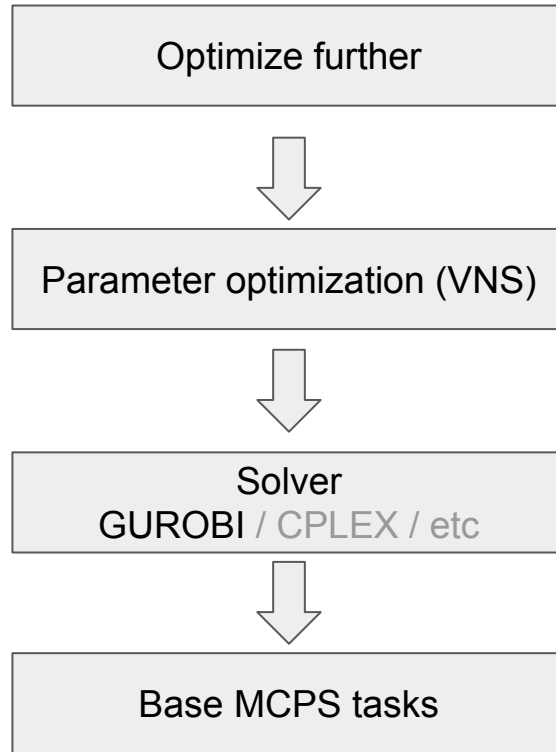
Let  $K = \{0, 1, 2, \dots, e\}$  be a set of event points, where  $e$  is a maximum event point, so we add the following variables:

1.  $t_{kc}$  — runtime of configuration  $c$  event point  $k$ .
2.  $x_{kc} = 1$  when and only when, configuration  $c$  is performed in event point  $k$ , 0 else. Let us assume that, zero configuration is performed in zero event point:  
 $x_{00} = 1$  и  $x_{0c} = 0$ ,  $c \in C$ .
3.  $y_{jk} = 1$  when and only when job  $j$  starts in event point  $k$ , 0 else.

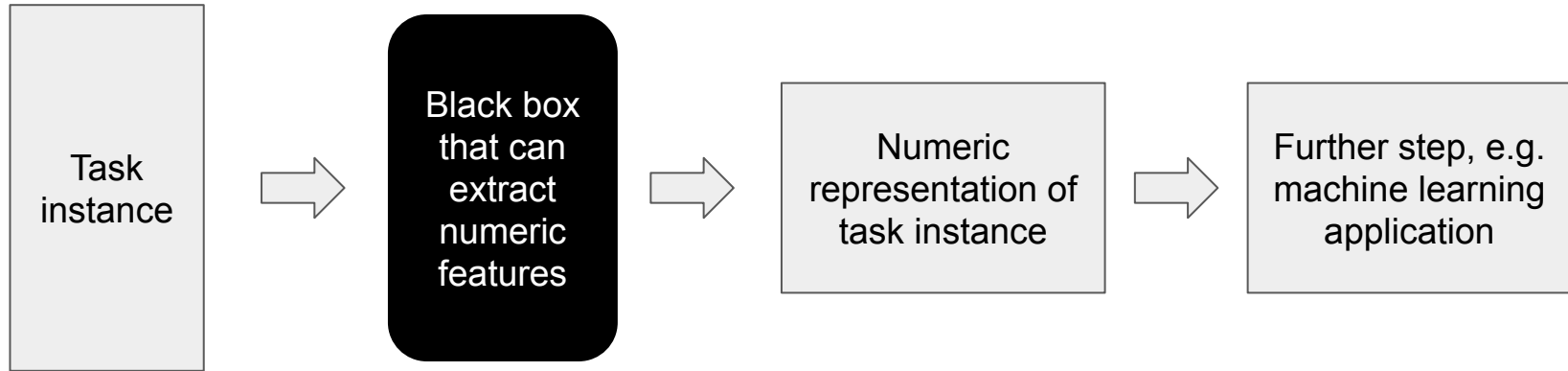
# Optimization process



# Optimization process



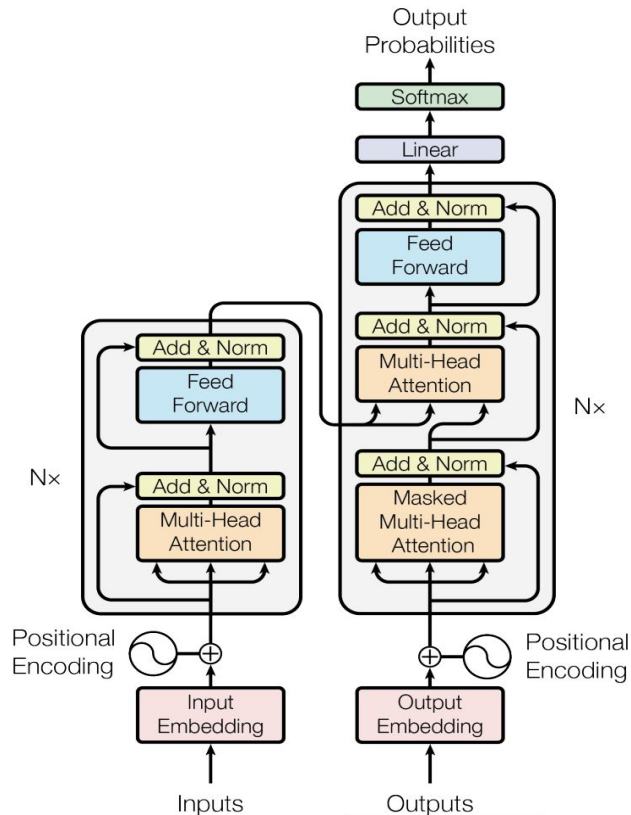
# Optimization process



# LLM (Large Language Models)

## BERT

Encoder

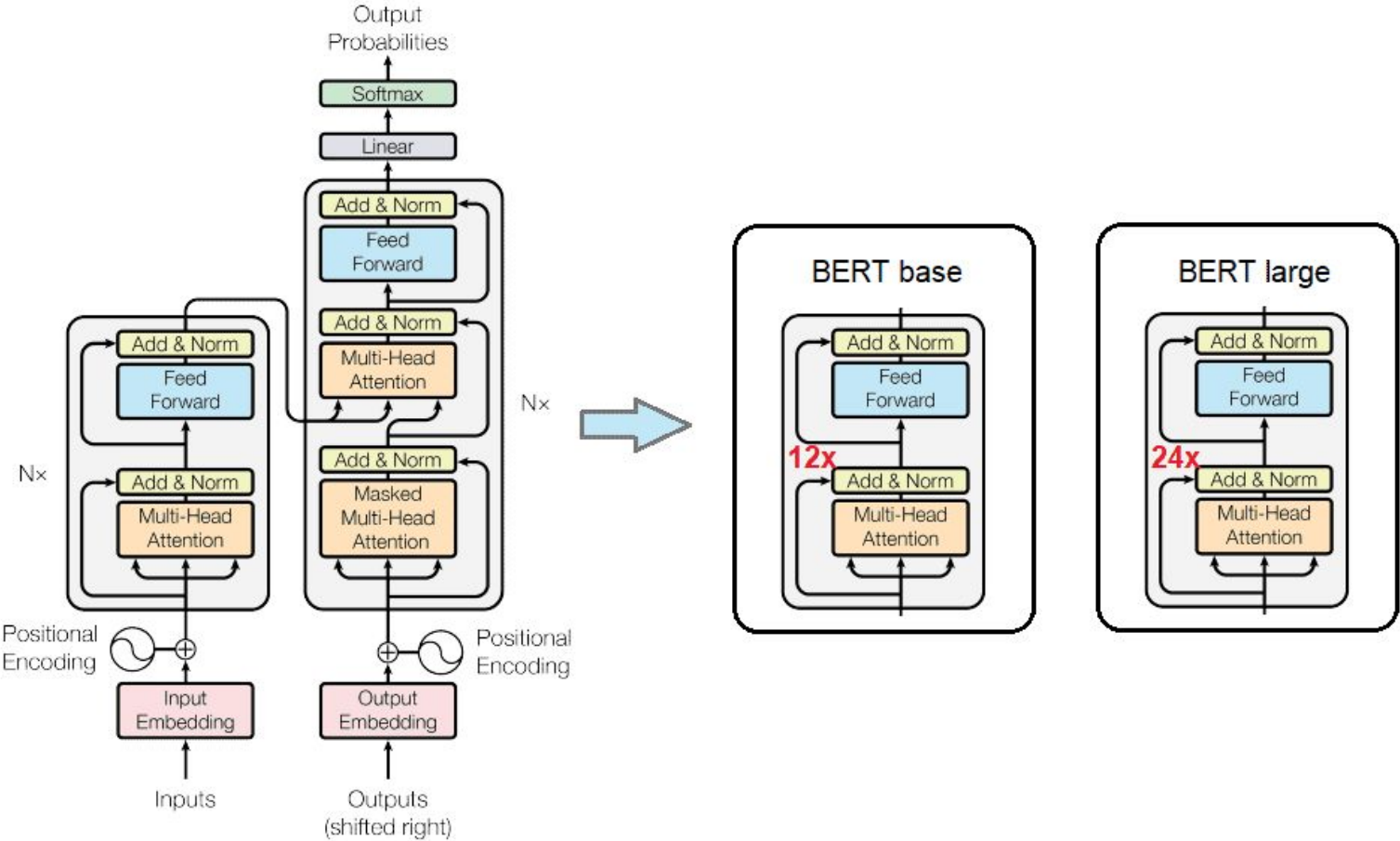


## GPT

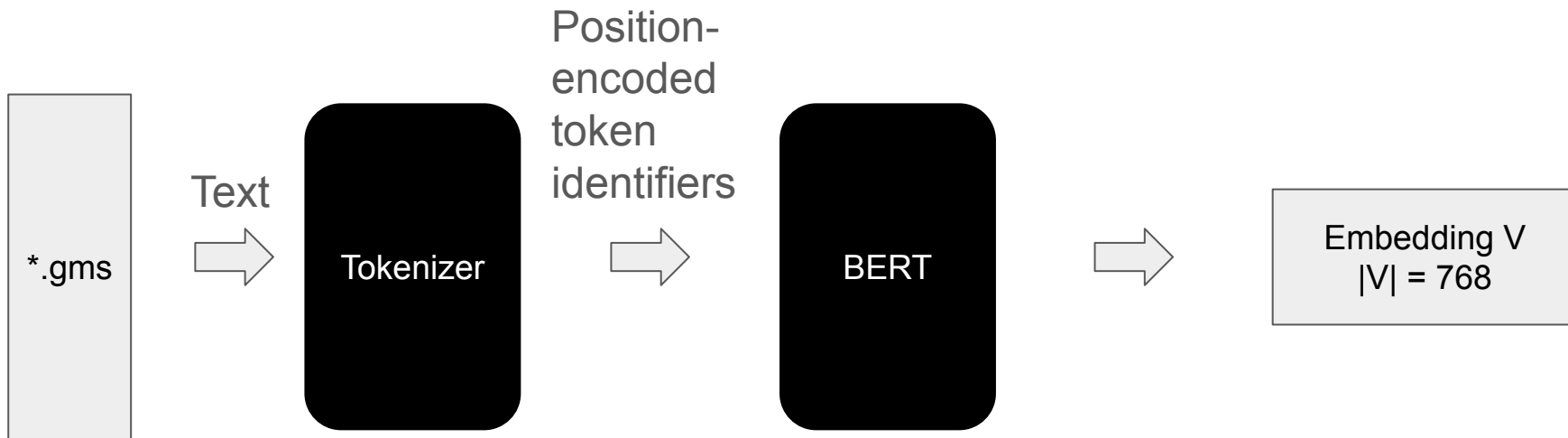
Decoder



# BERT(Bidirectional Encoder Representations from Transformers)



# BERT(Bidirectional Encoder Representations from Transformers)



# Where do I pass embeddings?

$P = (p_1, \dots, p_k)$  - solver's parameter vector (*configuration*)

Regression task:

$P_i^* = V_i W$ ,  $i = 1, \dots, J$ , where  $J$  is a number of different individual tasks

$V_i = \text{BERT}(\text{Tokenizer}(I_i))$

# Experimental data

Dimensionality	Optimization type		
	No	VNS	LLM
4 jobs	0.0701 s	0.0698 s	0.0695 s
7 jobs	117,7 s	–	117,04 s

Table 2. Average runtime

# Conclusions

- Since the process of obtaining embeddings and training linear regression was performed only for tasks with 4 jobs the unquestionable improvement was obtained only on this dimensionality;
- Bad generalization is seen on higher dimensionalities, probably because it's interpolation so it might be inaccurate;
- Assumption: for the purpose of increasing prediction quality it's best to “show” some amount of high-dimensional tasks, so the linear model can “extrapolate” the knowledge on in-between dimensionalities;
- Going for more parameters will make us face imminent curse of dimensionality, which is open problem in this case, since there are not so many task instances available;
- There is another way of obtaining vector representation, introduced in the paper about MIPLIB [3], those representations formed by looking inside of tasks during runtime and may be useful in our case.

# Thanks for your attention!

## References:

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