

# On analysis of models for customer order scheduling problem

Yu. Zakharova, A. Zakharov

Sobolev Institute of Mathematics SB RAS

The research was supported by Russian Science Foundation  
grant N 22-71-10015.

## Input data

$m$  is the number of customers

$n$  is the number of products

$p_{ij} \geq 0$  is the duration of producing product  $j$  for customer  $i$

$s_{jj'} \geq 0$  is the setup time from product  $j$  to product  $j'$

$s'_j$  is the initial setup to product  $j$

$d_i$  is the due date for customer  $i$

$q_i$  is the weight for customer  $i$

## Solution representation

We define operation as a pair of customer and product type  $(i, j)$ ,  $i = 1, \dots, m$ ,  $j = 1, \dots, n$ . Operations should be scheduled without preemptions. A set of feasible solution  $\Pi$  consists of permutations of operations  $\pi = \{(i, j), i = 1, \dots, m, j = 1, \dots, n\}$ .

## Makespan

$$\max_{i=1,\dots,m} C_i$$

## Weighted throughput

$$\sum_{i \in O(\pi)} q_i$$

$O(\pi)$  is the subset of customers, for which  $C_i \leq d_i$  in  $\pi$

## Total completion time

$$\sum_{i=1}^m C_i$$

## Maximum lateness

$$\max_{i=1,\dots,m} \{C_i - d_i\}$$

## Previous Research

Hazir, O., Gunalay, Y., Erel, E. (2008). Customer order scheduling problem: a comparative metaheuristics study. *The International Journal of Advanced Manufacturing Technology*, 37, 589-598.

Erel, E., Ghosh, J. B. (2007). Customer order scheduling on a single machine with family setup times: Complexity and algorithms. *Applied Mathematics and Computation*, 185(1), 11-18.

Cetinkaya, F. C., Yeloglu, P., Catmakas, H. A. (2021). Customer order scheduling with job-based processing on a single-machine to minimize the total completion time.

Kovalenko, Y. V., Zakharov, A. O. (2020). The Pareto set reduction in bicriteria customer order scheduling on a single machine with setup times. In *Journal of Physics: Conference Series* (Vol. 1546, No. 1, p. 012087). IOP Publishing.

$$x_{ijk} = \begin{cases} 1, & \text{if product } j \text{ of customer } i \text{ is placed in position } k, \\ 0 & \text{otherwise;} \end{cases}$$

$t_k^f \geq 0$  is the completion time of an operation in position  $k \in N$ ;

$T_i^f \geq 0$  is the completion time of the service of customer  $i \in M$ ;

$$z_i = \begin{cases} 1, & \text{if the order of customer } i \text{ is made on time } (T_i^f \leq d_i), \\ 0 & \text{otherwise.} \end{cases}$$

# Mathematical Programming Model 1 (constraints)

$$\sum_{k=1}^{|N|} x_{ijk} = 1, \quad i = 1, \dots, m, \quad j = 1, \dots, n, \quad (1)$$

$$\sum_{i=1}^m \sum_{j=1}^n x_{ijk} = 1, \quad k = 1, \dots, |N|, \quad (2)$$

$$t_1^f \geq \sum_{i=1}^m \sum_{j=1}^n x_{ij1} (p_{ij} + s'_j), \quad (3)$$

$$t_k^f \geq t_{k-1}^f + p_{ij} + \sum_{i'=1}^m \sum_{j'=1}^n x_{i'j',k-1} s_{j'j} - H(1 - x_{ijk}), \quad (4)$$

$$k = 2, \dots, |N|, \quad i = 1, \dots, m, \quad j = 1, \dots, n,$$

$$T_i^f \geq t_k^f - H(1 - x_{ijk}), \quad (5)$$

$$k = 1, \dots, |N|, \quad i = 1, \dots, m, \quad j = 1, \dots, n,$$

$$H(1 - z_i) \geq T_i^f - d_i, \quad i = 1, \dots, m, \quad (6)$$

$$T_i^f \geq 0, \quad t_k^f \geq 0, \quad x_{ijk} \in \{0, 1\}, \quad y_i \in \{0, 1\}, \quad (7)$$

$$i = 1, \dots, m, \quad j = 1, \dots, n, \quad k = 1, \dots, |N|.$$

$$x_{ijk} = \begin{cases} 1, & \text{if product } j \text{ of customer } i \text{ is placed in position } k, \\ 0 & \text{otherwise;} \end{cases}$$

$$b_{j'jk} = \begin{cases} 1, & \text{if arise a setup between products } j' \\ & \text{and } j \text{ in position } k = 2, \dots, |N|, \\ 0 & \text{otherwise;} \end{cases}$$

$$z_{ik} = \begin{cases} 1, & \text{if customer } i \text{ is in the system in position } k, \\ 0 & \text{otherwise;} \end{cases}$$

$$U_{ijk'i'} = \begin{cases} 1, & \text{if both } x_{ijk} = 1 \text{ and } z_{i'k} = 1, \\ 0 & \text{otherwise;} \end{cases}$$

$$V_{j'jk'i'} = \begin{cases} 1, & \text{if both } b_{j'jk} = 1 \text{ and } z_{i'k} = 1, \\ 0 & \text{otherwise.} \end{cases}$$

# Mathematical Programming Model 2 (constraints)

$$\sum_{k=1}^{|N|} x_{ijk} = 1, \quad i = 1, \dots, m, \quad j = 1, \dots, n, \quad (8)$$

$$\sum_{i=1}^m \sum_{j=1}^n x_{ijk} = 1, \quad k = 1, \dots, |N|, \quad (9)$$

$$kx_{ijk} \leq a_i, \quad i = 1, \dots, m, \quad j = 1, \dots, n, \quad k = 1, \dots, |N|, \quad (10)$$

$$\sum_{k=1}^{|N|} z_{ik} = a_i, \quad i = 1, \dots, m, \quad (11)$$

$$z_{i,k+1} \leq z_{ik}, \quad i = 1, \dots, m, \quad k = 1, \dots, |N| - 1, \quad (12)$$

$$\sum_{i=1}^m x_{ijk} + \sum_{i=1}^m x_{i,j',k-1} - 1 \leq b_{j'jk}, \quad j, j' = 1, \dots, n, \quad k = 1, \dots, |N|, \quad (13)$$



Operation is a pair  $(i, j)$ .

We denote the set of all operations as  $O$ , and let  $O_i$  corresponds to operations of customer  $i \in M$ . We also add two auxiliary operations 0 and  $nm + 1$  of zero duration.

$$x_{kl} = \begin{cases} 1, & \text{if operation } k \text{ immediately proceeds operation } l, \\ 0 & \text{otherwise;} \end{cases}$$

$\tilde{t}_k^f \geq 0$  is the completion time of operation  $k \in O$ ;

$T_i^f \geq 0$  is the completion time of servicing customer  $i \in M$ .

# Mathematical Programming Model 3 (constraints)

$$\sum_{l=1}^{nm+1} x_{kl} = 1, \quad k = 0, 1, \dots, nm, \quad (14)$$

$$\sum_{k=0}^{nm} x_{kl} = 1, \quad l = 1, \dots, nm + 1, \quad (15)$$

$$\tilde{t}_k^f \geq \tilde{t}_l^f + s_{lk} + p_k - H(1 - x_{lk}), \quad k, l = 0, \dots, nm + 1, \quad (16)$$

$$T_i^f \geq \tilde{t}_k^f, \quad k \in O_i, \quad i = 1, \dots, m, \quad (17)$$

$$H(1 - z_i) \geq T_i^f - d_i, \quad i = 1, \dots, m, \quad (18)$$

$$T_i^f \geq 0, \quad \tilde{t}_k^f \geq 0, \quad x_{kl} \in \{0, 1\}, \quad y_{kl} \in \{0, 1\}, \quad z_i \in \{0, 1\}, \quad (19)$$
$$i = 1, \dots, m, \quad k, l = 1, \dots, nm + 1.$$

$$x_{kl} = \begin{cases} 1, & \text{if operation } k \text{ immediately proceeds operation } l, \\ 0 & \text{otherwise;} \end{cases}$$

$$y_{kl} = \begin{cases} 1, & \text{if operation } k \text{ proceeds operation } l, \\ 0 & \text{otherwise.} \end{cases}$$

$$xy_{klq} = \begin{cases} 1, & \text{if } x_{kl} = 1 \text{ and } y_{lq} = 1, \\ 0 & \text{otherwise.} \end{cases}$$

# Mathematical Programming Model 4 (constraints)

$$\sum_{l=1}^{nm+1} x_{kl} = 1, \quad k = 0, 1, \dots, nm, \quad (20)$$

$$\sum_{k=0}^{nm} x_{kl} = 1, \quad l = 1, \dots, nm + 1, \quad (21)$$

$$\tilde{t}_q^f = \sum_{k,l=0}^{nm+1} xy_{klq} s_{kl} + \sum_{l=0}^{nm+1} y_{lq} p_l + p_q, \quad q = 0, \dots, nm + 1, \quad (22)$$

$$T_i^f \geq \tilde{t}_k^f, \quad k \in O_i, \quad i = 1, \dots, m, \quad (23)$$

$$xy_{klq} \geq y_{lq} + x_{kl} - 1, \quad k, l, q = 0, \dots, nm + 1, \quad (24)$$

$$\sum_{k=1}^{nm} y_{kq} - \sum_{k=1}^{nm} y_{kl} + (nm + 3)x_{ql} \leq nm + 2, \quad q, l = 1, \dots, nm, \quad (25)$$

$$y_{kq} + y_{qk} = 1, \quad q, k = 0, \dots, nm + 1, \quad q \neq k. \quad (26)$$

$$H(1 - z_i) \geq T_i^f - d_i, \quad i = 1, \dots, m, \quad (27)$$

$$T_i^f \geq 0, \quad \tilde{t}_k^f \geq 0, \quad x_{kl} \in \{0, 1\}, \quad y_{kl} \in \{0, 1\}, \quad z_i \in \{0, 1\}, \quad (28)$$

$$x_{kl} \leq y_{kl}, \quad k, l = 0, \dots, nm + 1, \quad (29)$$

$$x_{kq} + x_{qk} \leq 1, \quad k, q = 0, \dots, nm + 1, \quad (30)$$

$$\sum_{k=0}^{nm} y_{kq} + \sum_{k=1}^{nm+1} y_{qk} = nm + 1, \quad q = 1, \dots, nm. \quad (31)$$

# Special case (zero setup)

$$t_k^f = \sum_{l=0}^{nm} y_{lk} p_l + p_k, \quad k = 1, \dots, nm + 1, \quad (32)$$

$$y_{kq} + y_{ql} + y_{lk} \leq 2, \quad q, l, k = 1, \dots, nm. \quad (33)$$

Thank you for your attention!